NUMERICAL EVALUATION OF SPECIAL POWER SERIES INCLUDING THE NUMBERS OF LYNDON WORDS: AN APPROACH TO INTERPOLATION FUNCTIONS FOR APOSTOL-TYPE NUMBERS AND POLYNOMIALS

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Dedicated to Walter Gautschi on the occasion of his 90th birthday

Abstract. Because the Lyndon words and their numbers have practical applications in many different disciplines such as mathematics, probability, statistics, computer programming, algorithms, etc., it is known that not only mathematicians but also statisticians, computer programmers, and other scientists have studied them using different methods. Contrary to other studies, in this paper we use methods associated with zeta-type functions, which interpolate the family of Apostol-type numbers and polynomials of order $k$. Therefore, the main purpose of this paper is not only to give a special power series including the numbers of Lyndon words and binomial coefficients but also to construct new computational algorithms in order to simulate these series by numerical evaluations and plots. By using these algorithms, we provide novel computational methods to the area of combinatorics on words including Lyndon words. We also define new functions related to these power series, Lyndon words counting numbers, and the Apostol-type numbers and polynomials. Furthermore, we present some illustrations and observations on approximations of functions by rational functions associated with Apostol-type numbers that can provide ideas on the reduction of the algorithmic complexity of these algorithms.

Key words. Lyndon words, special numbers and polynomials, Apostol-type numbers and polynomials, arithmetical function, interpolation function, zeta type function

AMS subject classifications. 03D40, 05A15, 11A25, 11B68, 11B83, 11S40, 11Y16, 65Q30, 68R15

1. Introduction. The main motivation of this article is to construct not only new formulas for special power series representations involving the Lyndon words counting numbers and the Apostol-type numbers and polynomials but also algorithms for the numerical computations using these formulas. In addition, we define new functions including these infinite series and Lyndon words counting numbers, and their numerical values are computed and plotted with the help of constructed computational algorithms.

Let us start by giving a brief description of the Lyndon words, one of the main motivations of this article, and the related definitions of the generating functions of the numbers that count them. The $k$-ary Lyndon words of length $n$, represented by a primitive necklace consisting of $n$ beads of $k$ different colors, are the lexicographically smallest elements in the set derived from all primitive words having length $n$ over the $k$-letter alphabet. Here, primitive words means that a word cannot be written as a positive power of its subwords. For instance, let us consider the alphabet $\{0, 1\}$. All 2-ary Lyndon words of length 5 which are derived from this alphabet are given as follows: $\{00001, 00011, 00101, 00111, 01011, 01111\}$. It is clear that the elements of this finite set are primitive words (cf. [3, 4, 6, 7, 11] and the references cited therein).
In this paper, $L_k(n)$ denotes the numbers of $k$-ary Lyndon words of length $n$. These numbers are computed by the following formula (cf. [3, 4, 6, 7]):

\[
L_k(n) = \frac{1}{n} \sum_{d \mid n} \mu \left( \frac{n}{d} \right) k^d,
\]

where the sum is over all positive divisors of $n$ and $\mu$ is the Möbius function, which is an arithmetic function defined by (cf. [2])

\[
\mu(1) = 1,
\]

and for $n > 1$ if we write $n = p_1^{a_1} p_2^{a_2} \cdots p_k^{a_k}$, where $p_1, p_2, \ldots, p_k$ are $k$ distinct primes, then

\[
\mu(n) = \begin{cases} 
(-1)^k & \text{if } a_1 = a_2 = \cdots = a_k = 1, \\
0 & \text{if } n \text{ is the product of non-distinct primes (i.e., } n \text{ is not square-free).}
\end{cases}
\]

The paper is organized as follows. In Section 2 we give some auxiliary results. Two new special power series, including the numbers of Lyndon words and binomial coefficients, are defined in Section 3. We give new relations and identities related to these series including zeta-type functions, the Apostol-type numbers of order $k$, and the Apostol-type polynomials of order $k$. Section 4 is devoted to some algorithms for computing the values of the functions representing these special power series by using a modification of recurrence formulas for the family of Apostol-type numbers and polynomials of order $k$. Finally, in Section 5, by using our algorithms, we provide some numerical evaluations and plots. Furthermore, we present some illustrations and observations for the approximation by rational functions of functions representing special power series that can provide ideas on the reduction of the algorithmic complexity of our algorithms.

2. Some auxiliary results. We start this section with descriptions and formulas for the Apostol-type numbers and polynomials, including their generating functions and their interpolation functions.

The Apostol-type numbers and polynomials of order $k$ have been defined by the second author [14] as

\[
F_w(t; \lambda; k) = \frac{1}{(\lambda e^t + \lambda^{-1} e^{-t} + 2)^k} = \sum_{n=0}^{\infty} W_n^{(k)}(\lambda) \frac{t^n}{n!},
\]

and

\[
G_w(t, x, k; \lambda) = e^{tx} F_w(t; \lambda; k) = \sum_{n=0}^{\infty} W_n^{(k)}(x; \lambda) \frac{t^n}{n!},
\]

respectively, where $k$ is a non-negative integer and $\lambda$ is a real or complex parameter. A relation between the numbers $W_n^{(k)}(\lambda)$ and the polynomials $W_n^{(k)}(x; \lambda)$ is given in the following theorem [14]:
THEOREM 2.1. We have

\[ W_n^{(k)}(x; \lambda) = \sum_{m=0}^{n} \binom{n}{m} x^{n-m} W_m^{(k)}(\lambda). \]

Setting \( x = 0 \) in (2.1), we obtain

\[ W_n^{(k)}(\lambda) = W_n^{(k)}(0; \lambda). \]

A computational formula for the numbers \( W_n^{(k)}(\lambda) \) is given by the next theorem [15]:

THEOREM 2.2.

(2.2)

\[ W_n^{(c+d)}(\lambda) = \sum_{m=0}^{n} \binom{n}{m} W_m^{(c)}(\lambda) W_{n-m}^{(d)}(\lambda). \]

By using the above formula for the numbers \( W_n^{(k)}(\lambda) \), some numerical values of these numbers are obtained (cf. [15]):

\[ W_0^{(2)}(\lambda) = \frac{\lambda^2}{(\lambda + 1)^2}, \quad W_1^{(2)}(\lambda) = \frac{2\lambda^2 (1 - \lambda)}{(\lambda + 1)^3}, \quad W_2^{(2)}(\lambda) = \frac{4\lambda^2 (\lambda^2 - 3\lambda + 1)}{(\lambda + 1)^6}, \]

\[ W_0^{(3)}(\lambda) = \frac{\lambda^3}{(\lambda + 1)^6}, \quad W_1^{(3)}(\lambda) = \frac{3\lambda^3 (1 - \lambda)}{(\lambda + 1)^7}, \quad W_2^{(3)}(\lambda) = \frac{3\lambda^3 (3\lambda^2 - 8\lambda + 3)}{(\lambda + 1)^8}. \]

It should be noted that

\[ W_n(\lambda) = W_n^{(1)}(\lambda) \quad \text{and} \quad W_n(x; \lambda) = W_n^{(1)}(x; \lambda). \]

The next theorem states a recurrence relation for the numbers \( W_n(\lambda) \) (cf. [14]):

THEOREM 2.3. Let \( n \) be a positive integer, and let

\[ W_0(\lambda) = \frac{\lambda}{(\lambda + 1)^2}. \]

Then the recurrence relation

(2.3)

\[ 2W_n(\lambda) + \lambda \sum_{m=0}^{n} \binom{n}{m} W_m(\lambda) + \lambda^{-1} \sum_{m=0}^{n} (-1)^{n-m} \binom{n}{m} W_m(\lambda) = 0 \]

holds.

By (2.3), the following numerical values of the numbers \( W_n(\lambda) \) can be found:

\[ W_1(\lambda) = -\frac{\lambda (\lambda - 1)}{(\lambda + 1)^2}, \quad W_2(\lambda) = \frac{\lambda (\lambda^2 - 4\lambda + 1)}{(\lambda + 1)^4}, \]

\[ W_3(\lambda) = -\frac{\lambda (\lambda^3 - 11\lambda^2 + 11\lambda - 1)}{(\lambda + 1)^5}, \ldots \]

The well-known Apostol-Bernoulli polynomials \( B_n^{(k)}(x; \lambda) \) and the Apostol-Euler polynomials \( \mathcal{E}_n^{(k)}(x; \lambda) \) can be expressed in terms of the polynomials \( W_n^{(k)}(x; \lambda) \). Indeed, the following relations hold (cf. [8]):

\[ B_n^{(2k)}(x; \lambda) = \frac{(-1)^k (n + 2k)_{2k}}{\lambda^k} W_n^{(k)}(x - k; -\lambda), \]
\[ E_n^{2k}(x; \lambda) = \left( \frac{4}{\lambda} \right)^k W_n^{(k)}(x - k; \lambda). \]

Otherwise, these polynomials of order \( k \), \( B_n^{(k)}(x; \lambda) \) and \( E_n^{(k)}(x; \lambda) \), are defined by the generating functions

\[ \left( \frac{t}{\lambda e^t - 1} \right)^k e^{xt} = \sum_{n=0}^{\infty} B_n^{(k)}(x; \lambda) \frac{t^n}{n!} \]

and

\[ \left( \frac{2}{\lambda e^t + 1} \right)^k e^{xt} = \sum_{n=0}^{\infty} E_n^{(k)}(x; \lambda) \frac{t^n}{n!} \]

(cf. \([1, 5, 12, 16, 17, 18]\)).

**Remark 2.4.** For other properties, relations, and identities related to the numbers \( W_n(\lambda) \) and \( W_n^{(k)}(\lambda) \) as well as to the polynomials \( W_n(x; \lambda) \) and \( W_n^{(k)}(x; \lambda) \), the reader can consult \([8, 9, 14, 15]\).

In addition to the above preliminaries about the Apostol-type numbers and polynomials \( W_n^{(k)}(\lambda) \) and \( W_n^{(k)}(x; \lambda) \), we now give their interpolation functions

\[ \zeta_w(s, k; \lambda) = \sum_{m=0}^{\infty} (-1)^m \binom{m+2k-1}{m} \frac{\lambda^{m+k}}{(m+k)!} \]

and

\[ \zeta_w(s, x, k; \lambda) = \sum_{m=0}^{\infty} (-1)^m \binom{m+2k-1}{m} \frac{\lambda^{m+k}}{(x+m+k)!}, \]

respectively, where \( \lambda, s \in \mathbb{C} \) with \( |\lambda| < 1 \) and \( \Re(s) > 1 \) (cf. \([9]\)). Special cases of the above functions, related to the Hurwitz-Lerch zeta function, are given in \([9]\). For \( n \in \mathbb{N} \), we also have (cf. \([9]\))

\[ \zeta_w(-n, k; \lambda) = W_n^{(k)}(\lambda), \]

\[ \zeta_w(-n, x, k; \lambda) = W_n^{(k)}(x; \lambda). \]

### 3. Two special power series including the numbers of Lyndon words.

In this section, we define two new special power series involving the numbers of Lyndon words and binomial coefficients. We give relations between these series and zeta-type functions. In the next section, we give algorithms for computing the values of these series by using modifications of the corresponding recurrence formulas for the numbers \( W_n^{(k)}(\lambda) \).

First, we define a power series involving the numbers \( L_k(n) \) and the binomial coefficients.

**Definition 3.1.** Let \( \lambda \in \mathbb{C} \) with \( |\lambda| < 1 \), and let \( n \) be a positive integer. We define

\[ G(\lambda, n, k) = \sum_{m=0}^{\infty} \frac{(-2k)^m}{m!} L_{m+k}(n) \lambda^{m+k}. \]

Since

\[ \binom{-2k}{m} = (-1)^m \binom{m+2k-1}{m}, \]

we have

\[ G(\lambda, n, k) = \sum_{m=0}^{\infty} \frac{(-2k)^m}{m!} L_{m+k}(n) \lambda^{m+k}. \]
another representation of (3.1) is valid:

\[ G(\lambda, n, k) = \sum_{m=0}^{\infty} (-1)^m \binom{m + 2k - 1}{m} L_{m+k}(n) \lambda^{m+k}. \]

Combining (1.1) with (3.1) we get

\[ G(\lambda, n, k) = \frac{1}{n} \sum_{d|n} \mu \left( \frac{n}{d} \right) \sum_{m=0}^{\infty} \left( -2k \right) \binom{m}{m} \lambda^{m+k} (m+k)^d. \]

Using (2.4) in the above equation yields

\[ G(\lambda, n, k) = \frac{1}{n} \sum_{d|n} \mu \left( \frac{n}{d} \right) \zeta_w(-d, k; \lambda). \]

Finally, combining (2.6) with the above equation we obtain the following result:

**Theorem 3.2.** Let \( n \) be a positive integer. Then we have

\[ (3.2) \quad G(\lambda, n, k) = \frac{1}{n} \sum_{d|n} \mu \left( \frac{n}{d} \right) W_d^{(k)}(\lambda). \]

By substituting \( k = 1 \) into (3.2), we get the corollary:

**Corollary 3.3.**

\[ G(\lambda, n, 1) = \sum_{m=0}^{\infty} (-1)^m (m+1) L_{m+1}(n) \lambda^{m+1} \]

or

\[ G(\lambda, n, 1) = \frac{1}{n} \sum_{d|n} \mu \left( \frac{n}{d} \right) W_d(\lambda). \]

If \( n \) is replaced by a prime number \( p \), then (3.2) reduces to the following:

**Corollary 3.4.** Let \( p \) be a prime number. Then we get

\[ (3.3) \quad G(\lambda, p, k) = \frac{W_p^{(k)}(\lambda) - W_1^{(k)}(\lambda)}{p}. \]

**Remark 3.5.** Let \( k = 1 \) in (3.3). Then, for \( p = 2 \) and \( p = 3 \) it reduces to

\[ G(\lambda, 2, 1) = \frac{\lambda^2 (\lambda - 2)}{(\lambda + 1)^2}, \]

and

\[ G(\lambda, 3, 1) = \frac{4\lambda^2 (\lambda - 1)}{(\lambda + 1)^3}, \]

respectively.

Now, we define a new family of polynomials \( \mathcal{L}_n(x, m, k) \) of degree \( n \):

**Definition 3.6.** The polynomials \( \mathcal{L}_n(x, m, k) \) of degree \( n \) are defined by

\[ (3.4) \quad \mathcal{L}_n(x, m, k) = \frac{1}{n} \sum_{d|n} \mu \left( \frac{n}{d} \right) \sum_{j=0}^{d} \binom{d}{j} (m+k)^{d-j} x^j. \]
If \( n = p \) (\( p \) is a prime number), then (3.4) yields:

**Corollary 3.7.** Let \( p \) be a prime number. Then

\[
\mathcal{L}_p(x, m, k) = \frac{(x + m + k)^p - (x + m + k)}{p}.
\]

We also define a power series involving the polynomials \( \mathcal{L}_n(x, m, k) \):

**Definition 3.8.** Let \( \lambda \in \mathbb{C} \) with \( |\lambda| < 1 \), and let \( n \) be a positive integer. We define

\[
(3.5) \quad \mathcal{H}(x; \lambda, n, k) = \sum_{m=0}^{\infty} \left( -\frac{2k}{m} \right) \mathcal{L}_n(x, m, k) \lambda^{m+k}.
\]

Combining (3.5) with (3.4) yields

\[
\mathcal{H}(x; \lambda, n, k) = \frac{1}{n} \sum_{d|n} \mu \left( \frac{n}{d} \right) \sum_{m=0}^{\infty} \left( -\frac{2k}{m} \right) \lambda^{m+k} \left( x + m + k \right)^d.
\]

By using (2.5) in the above equality, we obtain

\[
\mathcal{H}(x; \lambda, n, k) = \frac{1}{n} \sum_{d|n} \mu \left( \frac{n}{d} \right) \zeta_{m}(-d, x, k; \lambda).
\]

Finally, with (2.7) the last equality yields the following result:

**Theorem 3.9.** We have

\[
(3.6) \quad \mathcal{H}(x; \lambda, n, k) = \frac{1}{n} \sum_{d|n} \mu \left( \frac{n}{d} \right) W_d^{(k)}(x; \lambda).
\]

**Remark 3.10.** Substituting \( k = 1 \) into (3.6) we have

\[
\mathcal{H}(x; \lambda, n, 1) = \frac{1}{n} \sum_{d|n} \mu \left( \frac{n}{d} \right) W_d(x; \lambda).
\]

In addition, for \( x = 0 \) it reduces to

\[
\mathcal{H}(0; \lambda, n, k) = \mathcal{G}(\lambda, n, k).
\]

By combining (3.5) and (3.6), we obtain a connection between the polynomials \( \mathcal{L}_n(x, m, k) \) and \( W_d^{(k)}(x; \lambda) \):

**Corollary 3.11.** We have

\[
\sum_{m=0}^{\infty} \left( -\frac{2k}{m} \right) \mathcal{L}_n(x, m, k) \lambda^{m+k} = \frac{1}{n} \sum_{d|n} \mu \left( \frac{n}{d} \right) W_d^{(k)}(x; \lambda).
\]

**4. A modification of the recurrence formula for the numbers** \( W_n^{(k)}(\lambda) \) **and algorithms.** In order to find the numerical values of the functions \( \mathcal{G}(\lambda, n, k) \), in this section we provide some computational algorithms for the numbers \( W_n(\lambda) \) and \( W_n^{(k)}(\lambda) \).

For \( n \geq 1 \), a modification of (2.3) is given by

\[
(4.1) \quad W_n(\lambda) = W_0(\lambda) \sum_{m=0}^{n-1} \left[ (-1)^{n-m+1} \lambda^{-1} - \lambda \right] \left( \frac{n}{m} \right) W_m(\lambda).
\]
By using (4.1), we state Algorithm 1 for computing the numbers $W_n(\lambda)$.

**Algorithm 1** Let $n$ be a nonnegative integer and $\lambda \in \mathbb{C}$. This algorithm will return the numbers $W_n(\lambda)$ recursively.

```plaintext
procedure W_APOSTOL_TYPE_NUM(n: nonnegative integer, \lambda)
    Begin
        Local variable m : positive integer
        if $n = 0$ then
            return $\lambda \power (\lambda + 1, 2)$
        else
            return $W_APOSTOL_TYPE_NUM(0, \lambda)$
            $\Rightarrow \sum((1/\lambda) \cdot \power(-1, n - m + 1) - \lambda) \cdot \Binomial_Coef(n, m)$
            $\Rightarrow W_APOSTOL_TYPE_NUM(n - m, \lambda, 1, n)$
        end if
    end procedure
```

We provide some numerical values of the numbers $W_n(\lambda)$ computed by Algorithm 1:

\[
W_1 \left(\frac{1}{4}\right) = 0.096, \quad W_2 \left(\frac{1}{4}\right) = 0.0064, \quad W_3 \left(\frac{1}{4}\right) = -0.08832, \\
W_4 \left(\frac{1}{4}\right) = -0.11648, \quad W_5 \left(\frac{1}{4}\right) = 0.059136, \quad W_6 \left(\frac{1}{4}\right) = 0.5126656.
\]

By setting $c = k - 1$ and $d = 1$ in (2.2), we also have the following recurrence relation for the numbers $W_n^{(k)}(\lambda)$:

\[
(4.2) \quad W_n^{(k)}(\lambda) = \sum_{m=0}^{n} \binom{n}{m} W_m^{(k-1)}(\lambda) W_{n-m}(\lambda).
\]

By using (4.2), we obtain Algorithm 2 for computing the numbers $W_n^{(k)}(\lambda)$.

**Algorithm 2** Let $n$ be a nonnegative integer, $\lambda \in \mathbb{C}$, and $k$ be a positive integer. This algorithm will return the numbers $W_n^{(k)}(\lambda)$ recursively with the help of the procedure $\text{W}_APOSTOL_TYPE_NUM$ given by Algorithm 1.

```plaintext
procedure HIGHER_W_APOSTOL_TYPE_NUM(n: nonnegative integer, \lambda, k: positive integer)
    Begin
        Local variable m : nonnegative integer
        if $k = 1$ then
            return $W_APOSTOL_TYPE_NUM(n, \lambda)$
        else
            return $\sum(\Binomial_Coef(n, m) \cdot \text{HIGHER_W_APOSTOL_TYPE_NUM}(m, \lambda, k - 1))$
            $\Rightarrow \sum(\text{HIGHER_W_APOSTOL_TYPE_NUM}(n - m, \lambda, 1, m, 0, n))$
        end if
    end procedure
```
Some numerical values of the numbers $W_n^{(k)}(\lambda)$ computed by Algorithm 2 are as follows:

\[
W_0^{(2)} \left( \frac{1}{2} \right) = 0.0493827160494, \\
W_0^{(3)} \left( \frac{1}{3} \right) = 0.006591796875, \\
W_1^{(2)} \left( \frac{1}{2} \right) = 0.0329218106996.
\]

By using (3.2) related to the numbers $W_n^{(k)}(\lambda)$ and the Möbius function $\mu(n)$, we can also state Algorithm 3 for computing the values of the functions $G(\lambda, n, k)$.

**Algorithm 3** Let $\lambda \in \mathbb{C}$ with $|\lambda| < 1$, $n$ be a nonnegative integer, and $k$ be a positive integer. This algorithm will return $G(\lambda, n, k)$ given by (3.2) with the help of the procedure HIGHER_W_APOSTOL_TYPE_NUM given by Algorithm 2 and the Möbius function denoted as the procedure Mobius_Func.

```
procedure G_LYNDON_FUNC(\lambda, n: nonnegative integer, k: positive integer)
    Begin
        Local variable G ← 0
        for all positive divisors d of n do
            G ← G + Mobius_Func(n/d) * HIGHER_W_APOSTOL_TYPE_NUM(d, \lambda, k)
        end for
        return G
    end procedure
```

A computation by Algorithm 3 yields some values for the functions $G(\lambda, n, k)$:

\[
G \left( \frac{1}{2}, 2, 1 \right) = -0.0740740740741, \\
G \left( \frac{1}{3}, 3, 1 \right) = -0.0658436213992, \\
G \left( \frac{1}{2}, 5, 1 \right) = 0.0658436213992.
\]

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\]

5. **Numerical evaluation of the numbers $W_n(\lambda)$ and $W_n^{(k)}(\lambda)$ and the functions $G(\lambda, n, k)$**. In this section, by using our algorithms, we not only provide plots for the numbers $W_n(\lambda)$ and the functions $G(\lambda, n, k)$, but also give some numerical evaluations related to the numbers $W_n(\lambda)$, the numbers $W_n^{(k)}(\lambda)$, and the functions $G(\lambda, n, k)$. Moreover, we present some illustrations and observations on the approximation of the functions $G(\lambda, p, 1)$ by rational functions that can provide ideas for a reduction of the algorithmic complexity of our algorithms.

We assume that $\lambda$ is a real number. Firstly, by using Algorithm 1, we plot the numbers $W_n(\lambda)$ in the cases of $n = 0, 1, \ldots, 6$ and for varying $\lambda \in \left[ \frac{1}{10^n}, \frac{1}{10^{n+1}} \right]$. It is clear that the $W_n(\lambda)$ are rational functions of the variable $\lambda$. Thus, the curves in Figure 5.1 provide some information for analysing some characteristics of the rational functions $W_n(\lambda)$. 
Graphs of the numbers $W_n(\lambda)$ for $n = 0, 1, \ldots, 6$ and varying $\lambda \in [\frac{1}{100}, \frac{1}{10}]$.

Graphs of the functions $G(\lambda, p, 1)$ for the prime numbers $p = 2, 3, 5, 7$ and $\lambda \in [0, 5]$ obtained by Algorithm 3 are presented in Figure 5.2. This figure demonstrates the effects of the prime numbers on the shape of the curve for the selected range $\lambda \in [0, 5]$, and these curves provide information for analysing some characteristics of the functions $G(\lambda, p, 1)$.

Now, we provide an approximation to the functions $G(\lambda, p, 1)$ by rational functions with small errors to give ideas for a reduction of the algorithmic complexity of Algorithm 3 for prime numbers.
Let us assume that $|\lambda| < 1$, and let $p$ be a prime number. It is clear that the numbers $W_n(\lambda)$ are rational functions of the variable $\lambda$. Therefore, by setting

$$G_A(\lambda, p) = \frac{W_p(\lambda)}{p},$$

the following inequality holds true:

$$|G(\lambda, p, 1) - G_A(\lambda, p)| \leq \varepsilon_p,$$

where $\varepsilon_p = 1/p$.

We present some plots and numerical experiments in order to illustrate the approximations of the functions $G(\lambda, p, 1)$ by the rational functions $G_A(\lambda, p)$ with an error less than $\varepsilon_p = 1/p$. The rational functions $G_A(\lambda, p)$ for four different prime numbers ($p = 3, 5, 7, 11$) are presented by red lines in Figure 5.3. According to (5.1) for sufficiently large prime numbers $p$, $\varepsilon_p$ tends to zero, and the curves of the functions $G(\lambda, p, 1)$ and $G_A(\lambda, p)$ tend to overlap. This indicates that by using the rational functions $G_A(\lambda, p)$ instead of $G(\lambda, p, 1)$, Algorithm 3 can operate more efficiently for sufficiently large prime numbers.

\begin{figure}[h]
\centering
\includegraphics[width=\textwidth]{figure5.3.png}
\caption{Approximations of the function $G(\lambda, p, 1)$ by the rational function $G_A(\lambda, p)$ in the cases: (a) $p = 3$, $\varepsilon_3 = 1/3$; (b) $p = 5$, $\varepsilon_5 = 1/5$; (c) $p = 7$, $\varepsilon_7 = 1/7$; (d) $p = 11$, $\varepsilon_{11} = 1/11$.}
\end{figure}

**Remark 5.1.** In [10], Kucukoglu et al. constructed generating functions for the $k$-ary Lyndon words having prime number length with the help of the Apostol-Bernoulli numbers and other special numbers. Moreover, they gave an approximation to these generating functions by rational functions of the Apostol-Bernoulli numbers. In this paper, we give approximations
for functions representing our special power series by rational functions associated with Apostol-type numbers by using similar techniques as the ones in [10].

REFERENCES